## 6.4 Practice Questions: Short Truth Table Method

Test the following arguments for validity using the short truth table method:

1. 1. (P  $\rightarrow$  Q) 2. –Q 3. *∴ –*P

> Conclusion Premise 2 Premise 1

Р	Q	-P	-Q	P→Q
Т	Т	F	F	Т

This argument is valid. If the conclusion is false while Premise 1 is true, then Premise 2 must be false on pain of contradiction. Alternatively, if the conclusion is false and Premise 2 is true, then Premise 1 must be false.

- 2. 1. (P & Q)
  - 2. Q  $\rightarrow$  R
  - 3. ∴ R

Conclusion Premise 1 Premise 2

Р	Q	R)	P & Q	$Q \rightarrow R$
Т	т	F	Т	F

This argument is valid. First, we make the conclusion (R) false. Next, we set Premise 1 to true because it is a conjunction. Knowing that Premise 1 is true allows us to make P and Q true. Last, we try to make Premise 2 true. If Q is true and R is false, then  $Q \rightarrow R$  must be false as well. We have already made Q true while R is false, so we cannot make Premise 2 true without contradicting ourselves.

3. 1. (P v Q) 2. P  $\rightarrow$  R 3.  $\therefore$  Q

	Conclusion	Premise 1	Premise 2	
Р	Q	R	ΡvQ	$P \rightarrow R$
Т	F	Т	Т	Т

This argument is invalid. If Premise 1 is true, and Q is false, then P must be true. We may set Premise 2 to true, which results in R being true, and we have no contradiction with Q being false.

1. −(P v Q)
 2. −Q v R
 3. ∴ −R

				Premise 1	Conclusion	Premise 2	
Р	Q	R	ΡvQ	-(P v Q)	-Q	-R	–Q v R
F	F	Т	F	Т	Т	F	Т

This argument is invalid. First we make the conclusion -R false, which makes R true. Next, we make Premise 1 true because it is a negated disjunction. When an inclusive disjunction is negated, both of its conjuncts are false. Because Q is false, -Q is true. Finally, the second premise ( $-Q \vee R$ ) can be true: we already know that -Q is true, so the disjunction is true, whether or not R is true. Because both premises can be true while the conclusion is false, this argument is invalid 5. 1. (P & Q) → R

2	D	$\sim$
Z	(–P v –	·U)

3. ∴ R

					Premise 1		Premise 2	
Р	Q	R	-Р	-Q	(P & Q)	(P & Q) → R	(–P v –Q)	-(-P v -Q)
Т	Т	F	F	F	Т	F	F	Т

This argument is valid. First, we set the conclusion to false. Next, we set Premise 2 to true; because it is a negated inclusive disjunction, we know that both disjuncts must be alse, which results in P and Q both being true. Because both P and Q are true, (P & Q) is also true. Because (P & Q) is true and R is false, Premise 1 must be set to false.

6. 1. (P v Q) & (R v S)

2. P → -S

3. –Q

4. ∴ R

		Conclusion		Premise 3				Premise 1	Premise 2
Р	Q	R	S	-Q	<b>-</b> S	ΡvQ	R v S	(P v Q) & (R v S)	$P \rightarrow -S$ )
Т	F	F	Т	Т	F	Т	Т	Т	F

This argument is valid. First we make the conclusion false. Next we make Premise 3 true, which forces Q to be false. Next we make Premise 1 true, which means both (P v Q) and (R v S) must be true. Because R is false, S must be true for (R v S) to be true. Because S must be true, –S must be false. Next, because Q is false, P must be true for (P v Q) to be true. Finally, we try to make Premise 2 true. However, we are forced to set Premise 2 to false: we already determined that P has to be true and –S has to be false when setting Premise 1 to true. Because the antecedent is true while the consequent is false, P remise 2 must be set to false. As a result, because we cannot make all of the premises true while the conclusion is false, this argument is valid.

7. 1. (P v Q) & R 2. R  $\rightarrow -Q$ 3.  $\therefore$  P

Conclusion

Premise 1 Premise 2

Premise 4

Conclusion

Р	Q	R	-Q	(P v Q)	(P v Q) & R	R → –Q
F	F	Т	Т	F	F	Т

This argument is valid. First, we set the conclusion P to false. Next we try to make both premises true. If we make Premise 1 true, then R must be true. If R is true, and Premise 2 is true, then Q must be false. If P and Q are both false, then (P v Q) cannot be true. If (P v Q) cannot be true, then we cannot make Premise 1 true. So, at least one premise must be false when the conclusion is false.

8. 1. ( $P \lor Q$ ) 2.  $Q \rightarrow R$ 3.  $R \rightarrow S$ 4. -P5.  $\therefore$  ( $S \lor T$ )

Р	Q	R	S	Т	ΡvQ	$Q \rightarrow R$	$R \rightarrow S$	-P	(S v T)
F	Т	Т	F	F	Т	Т	F	Т	F

This argument is valid. First, we, set the conclusion (S v T) to false. For (S v T) to be false, both S and T must be false. Next, we try to make each premise true. For -P to be true, P must be false. For (P v Q) to be true when P is false, Q must be true. Because Q is true, and Premise 2 is true, R must be true. However, if R is true and S is false, Premise 3 must be false.

Premise 1 Premise 2 Premise 3

9. 1. (P & Q) → R
2. (P v S)
3. -S
4. (Q & T)
5. ∴ R
Conclusion

Premise 1	Premise 2	Premise 3	Premise 4

Р	Q	R	S	Т	(P & Q)	(P & Q) → R	(P v S)	-S	(Q & T)
Т	Т	F	F	Т	Т	F	Т	Т	Т

This argument is valid. First, we set the conclusion R to false. Next, we try to make all the premises true. We begin with the premises that will force us to assign values to simple statements. For Premise 3 to be true, S must be false. For Premise 4 to be true, we must set Q and T to true. For Premise 2 to be true while S is false, P must be true. From what we have filled in, we can set (P & Q) to true because both P and Q are true. Finally, Premise 1 must be false: R is false while (P & Q) is true. Because at least one premise must be false when the conclusion is false, this argument is valid.

- 10. 1.-P
  - 2. Q → R
  - 3. (P v Q)
  - 4. ∴ Q & R

Premise 1 Premise 2 Premise 3 Conclusion

Р	Q	R	-P	$Q \rightarrow R$	ΡvQ	Q & R
F	Т	F	Т	F	Т	F
F	F	F	Т	Т	F	F
F	F	Т	Т	Т	F	F

This argument is valid. We constructed three rows to explore the three possible ways (Q & R) can be false. In each instance, at least one premise must be false when (Q & R) is false, so this argument is valid.